1 METHOD AND APPARATUS FOR SERVING CONTENT FROM A 2 SEMI-TRUSTED SERVER

FIELD OF THE INVENTION

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- 4 This invention is directed to the field of computer
- 5 networks. It is more particularly directed to the
- 6 Internet and specifically to web-servers which store
- 7 and transmit information to their clients.

BACKGROUND OF THE INVENTION

Much of the communication in the Internet involves at least one client making a request to a web-server, and the web-server responding to the client's request. By web-server, we mean a device or collection of devices (e.g., datastore, directory, machines, and software) that communicates with a client using the HTTP Protocol. For a particular web application, we define an origin web-server to be a web-server that is completely trusted with the functions and data of a web application with regard to the security policy for the application.

As a way to shorten the length of time that the client must wait for a response and to lighten the load on the Internet and origin web-server, techniques have been developed to allow the client to be serviced by a proxy web-server, also referred to simply as a proxy, where

- 1 the proxy web-server is usually closer to the client or
- 2 more lightly loaded than the origin web-server.
- 3 Proxy web-servers can be integrated into the Internet
- 4 communication process in several different ways. In
- 5 some configurations, clients always make requests to a
- 6 proxy web-server rather than the origin web-server.
- 7 The proxy web-server may respond to the client,
- 8 fetching content from the origin web-server as
- 9 necessary, or the proxy web-server may refer the client
- 10 to another proxy web-server or the origin server if the
- $_{i=1}$ 1 proxy web-server is unable to satisfy the client's
- \square request. In other configurations, a client first makes
- a request to the origin web-server. The origin
- web-server may refer the client to a proxy for the
- current request or all future requests, or the origin
- web-server may respond to part of the client's request,
- but refer the client to a semi-trusted web-server for a
- 18 portion of the response.
- In most cases, the content offloaded to a proxy
- 20 web-server has been limited to non-sensitive data, so
- that access control schemes are not required.
- Non-sensitive data is defined as data which does not
- 23 require any access control, and may be accessible to
- 24 any user on the network. On a typical web-page,
- embedded images are an example of non-sensitive data.
- 26 On the other hand, restricted data or sensitive data is
- 27 defined as data which has some restrictions on who can
- obtain it. Examples of restricted data include pages
- that are obtained by subscription to a set of
- registered users, images that are available to a

- 1 restricted set of users, or data can is personalized
- 2 for a specific user.
- 3 Common subscription services and personalized content
- 4 on the Internet are increasing, and they should also
- 5 benefit from the performance gains afforded by proxy
- 6 web-servers. The restricted information requires the
- 7 proxy web-servers to have access control methods in
- 8 place, but the situation is complicated because in many
- 9 cases the proxy web-servers are not under the control
- of the content providers. Such proxy web-servers fall
- $_{\mathbb{R}}$ 11 into the class of semi-trusted web-servers. For a
- 42 particular web application, we define a semi-trusted
- web-server to be a web-server that is partially trusted
- 4 for the functions of the application with regard to the
- security policy for the application. In particular, a
- ₹16 semi-trusted web-server may be trusted for
- authorization, access to user identifiers, SSL
- 18 tunneling to an origin web-server for content, and
- non-sensitive transactions, but the semi-trusted
- $\frac{\mathbb{Z}_2^{0}}{\mathbb{Z}_2^{0}}$ web-server may not be trusted with long-term sensitive
- data such as user passwords or secret keys for an
- origin web-server.

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SUMMARY OF THE INVENTION

- It is therefore an aspect of the present invention to
- 25 provide a method to enable a client to access
- restricted information from an origin web-server
- through a semi-trusted web-server.

- 1 It is a further aspect of the invention to provide an
- 2 apparatus to enable a client to access restricted
- 3 information from an origin web-server through a
- 4 semi-trusted web-server.

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- 5 It is a further aspect of the invention to reduce the
- 6 security risk of using cookies for authentication.
- 7 It is a further aspect of the invention to protect data
- 8 stored at a semi-trusted web-server.
 - In an example of the invention, a client would use the method disclosed herein to gain access to restricted information at a semi-trusted web-server.

BRIEF DESCRIPTION OF THE DRAWINGS

- These and other aspects, features, and advantages of the present invention will become apparent upon further consideration of the following detailed description of the invention when read in conjunction with the drawing figures, in which:
- 18 Fig. 1 shows an example of an environment having a
- 19 client, an origin web-server, and a semi-trusted
- 20 web-server, and the problems associated with a client
- 21 accessing restricted information through the
- 22 semi-trusted web-server.
- 23 Fig. 2 shows the deployment of the method described in
- this invention for enabling a client to access
- 25 restricted information in the environment described by

- 1 Fig. 1. It shows the authenticator, creator,
- presentator, and correlator.
- Fig. 3 shows a first embodiment of this invention
- 4 illustrating a technique for presenting a client
- 5 credential for the case when the semi-trusted
- 6 web-server and the origin web-server are located in the
- 7 same name space domain, (e.g. proxy.company.com and
- 8 company.com respectively).
- 9 Fig. 4 shows a flowchart that illustrates the actions
- taken by the origin web-server for the embodiment shown
- 411 in Fig. 3.

- 12 Fig. 5 shows a flowchart that illustrates the actions
- taken by the semi-trusted web-server for the embodiment
- $\frac{1}{2}$ 14 shown in Fig. 3.
- Fig. 6 shows a second embodiment in which the client
- [16] credential is presented to the semi-trusted web-server
- 17 by the origin web-server and the origin server
 - 18 redirects the client to the semi-trusted web-server
 - 19 with HTTP redirection. This embodiment addresses the
 - case when the origin web-server and semi-trusted
 - 21 web-server are not located in the same domain.
 - 22 Fig. 7 shows a third embodiment in which correlation is
 - 23 aided by the origin web-server installing a first
 - 24 client-side program in the client browser that sends
 - 25 client-specific information to the origin web-server
 - 26 and by the semi-trusted web-server installing a second
 - 27 client-side program in the client browser that sends

- 1 client-specific information to the semi-trusted
- web-server.
- Fig. 8 shows the sequence of events for the embodiment
- 4 shown in Fig. 7.
- 5 Fig. 9 shows a flowchart that illustrates the actions
- 6 taken by the origin web-server for the embodiment shown
- 7 in Fig. 7.

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- 8 Fig. 10 shows a flowchart that illustrates the actions
- 9 taken by the semi-trusted web-server for the embodiment
- 410 shown in Fig. 7.
- Fig. 11 shows a flowchart that illustrates the actions
- taken by the origin web-server for creating a cookie
- that will be presented to the semi-trusted web-server.
- Fig. 12 shows a flowchart that illustrates the actions
- taken by the origin or semi-trusted web-server when
- 16 validating a cookie as part of the correlation
 - 17 procedures.
 - 18 Fig. 13 shows a flowchart that illustrates the actions
 - taken by the origin or semi-trusted web-server when
 - 20 updating a cookie as part of the procedures described
 - above.
 - 22 Fig. 14 illustrates the use of shared keys by
 - semi-trusted and origin web-servers.

- 1 Fig. 15 shows a flowchart that illustrates the actions
- taken by the origin web-server when creating secure
- 3 content for distribution through a semi-trusted
- 4 web-server.

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- 5 Fig. 16 shows a flowchart that illustrates the actions
- taken by the semi-trusted web-server when it serves
- 7 partly secure content.
- 8 Fig. 17 shows a flowchart that illustrates the actions
- 9 taken by the client when receiving secure content.
 - Fig. 18 illustrates the use of a new content-type that distinguishes secure content.
- Other objects and a better understanding of the invention may be realized by referring to the detailed description.

DESCRIPTION OF THE INVENTION

- 16 The present invention enables a client to access
- 17 restricted information through a semi-trusted
- 18 web-server. A typical environment in which the access
- 19 occurs is illustrated in Fig. 1. The figure shows a
- client, an origin web-server, and a semi-trusted
- 21 web-server which are connected to a core network such
- 22 as the Internet.
- The client 101, semi-trusted web-server 104, and origin
- web-server 103 are connected to a core network 102.

The client, semi-trusted web-server, and origin
web-server can be directly connected to the core
network as exemplified in the figure, or they can be
connected via intermediary firewalls, routers, and
subnetworks. The semi-trusted web-server may be on the
route that packets in the network follow when the
client communicates with the origin web-server.

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To shorten the length of time that the client must wait for a response to requests and to lighten the load on the origin web-server, the semi-trusted web-server may service the requests of the client. Generally, a semi-trusted web-server is chosen to service the requests of a client if the semi-trusted web-server is closer to the client than the origin web-server or if the semi-trusted web-server is less loaded than the origin web-server. In some configurations, clients always make requests to a semi-trusted web-server rather than the origin web-server. The semi-trusted web-server may respond to the client, fetching content from the origin web-server as necessary, or the semi-trusted web-server may refer the client to another semi-trusted web-server or the origin server if the semi-trusted web-server is unable to satisfy the client's request. In other configurations, a client first makes a request to the origin web-server. origin web-server may refer the client to a semi-trusted web-server for the current request or all future requests, or the origin web-server may respond to part of the client's request but refer the client to a semi-trusted web-server for a portion of the response.

In Fig. 1, the semi-trusted web-server is not necessarily considered to be under the control of the content provider or the origin web-server such that the semi-trusted web-server is not trusted with regard to the security policy for long-term sensitive data. However, the semi-trusted web-server is trusted for authorization, user identifiers, non-sensitive transactions, and SSL tunneling for content. Because the semi-trusted web-server is semi-trusted by the origin web-server, access control methods are needed to protect the restricted information at the semi-trusted web-server and to protect sensitive user credentials. Our patent provides techniques to provide access control to restricted information at the semi-trusted web-server without risking long-term sensitive data at the semi-trusted web-server such as user passwords.

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To enable a client to access restricted information at the semi-trusted web-server without risking long-term sensitive data, additional functions are needed. Fig. 2 shows the structure of any machine that is running in the network that can deploy the invention disclosed herein. The machine 201 may be distributed among the client, semi-trusted web-server, origin web-server, trusted agents of the client, trusted agents of the semi-trusted web-server, and trusted agents of the origin web-server. Any device 201 implementing this invention consists of four components including an authenticator 202, credential creator 203, credential presentator 204, and credential correlator 205. In some embodiments of this invention, the authenticator

202 may be provided by external services which are 1 2 already present in the network. The authenticator 202 3 first authenticates the client 101 to the origin web-server 103. The credential creator 203 generates 4 5 credential for the client to be used for subsequent communications. The credential presenter 204 6 7 communicates the client credential to the semi-trusted 8 web-server 104. The credential correlator 205 9 correlates the client credential with the accessing

10 client and the client user identifier.

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When a client requests restricted information, the client is referred to the origin web-server for authentication. After credential creation, the client is referred back to the semi-trusted web-server to repeat the request and the credential presented to the semi-trusted web-server. After the credential correlation, the semi-trusted web-server responds to the client's request by providing the requested content if the credential is valid.

Fig. 3 shows a first embodiment of this invention illustrating a technique for presenting a client credential for the case when the semi-trusted web-server 300 and the origin web-server 310 are located in the same name space domain (e.g. proxy.company.com and company.com respectively). In this embodiment, the client makes a request to the origin web-server. The request is sometimes referred to the origin web-server by the semi-trusted web-server. The client is authenticated to the origin web-server with the authentication 320 component.

1 the origin web-server, the credential creation 325 2 component creates a client credential in the form of an 3 HTTP cookie which includes client-specific environment 4 information such as the apparent client Internet 5 Protocol (IP) address and the HTTP header information. 6 The origin web-server sets the cookie on the client and 7 the client is referred to the semi-trusted web-server by the origin web-server using HTTP redirection. 8 9 Because the semi-trusted web-server and origin 10 web-server are in the same name-space domain, the 11 credential cookie will be presented to the semi-trusted 12 web-server 315 when the client makes a request to the **13** semi-trusted web-server. At the semi-trusted 1114 web-server, the credential correlation 330 component 1115 correlates the presented cookie to the client-specific 16 environment information such as the apparent client IP **#17** address and the HTTP header information. If the cookie 18 is valid, the semi-trusted web-server responds to the **1119** request using the access control of the user identifier 20 specified in the cookie. The format and validation of 21 the cookie are illustrated in Fig. 11.

Fig. 4 shows a flowchart that illustrates the actions taken by the origin web-server 310 for the embodiment shown in Fig. 3. The flowchart is entered in the step 401 whenever the device implementing the embodiment is started at the origin web-server 310. In step 405, the origin web-server waits for messages from a client. Upon receiving a message, in step 410 the origin web-server checks to see if the client is authenticated. The client may be authenticated by presenting a valid client credential as illustrated in

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1 Fig. 12, or the client may be authenticated by another 2 scheme. If the client is not authenticated in step 410, 3 then in step 425 the origin web-server initiates 4 authentication and returns to step 410. If the client 5 is authenticated in step 410, then step 415 is 6 executed. In step 415, the origin web-server uses 7 client-specific environment information to create a 8 client credential which is stored in a cookie that is 9 set on the client. Next, in step 420, the origin 10 web-server refers the client to a semi-trusted 11 web-server. Finally, the origin web-server returns to 12 step 405 and waits for another message.

> Fig. 5 shows a flowchart that illustrates the actions taken by the semi-trusted web-server 300 for the embodiment shown in Fig. 3. The flowchart is entered in step 500 whenever the device implementing the embodiment is started at the semi-trusted web-server In step 505, the semi-trusted web-server waits for messages from clients. Upon receiving a message, the semi-trusted web-server checks to see if the client has submitted a cookie containing a valid client credential in step 510. Step 510 involves correlating the client credential in the cookie with client-specific information that the semi-trusted web-server obtains from the client as illustrated in Fig. 12. If the client credential is not valid in step 510, then step 530 is executed. In step 530, the semi-trusted web-server refers the client to the origin web-server and returns to step 505. If the client credential is valid in step 510, then step 515 is executed. In step 515, the cookie and client

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credential is updated as illustrated in Fig. 13. After 1 2 step 515, the semi-trusted web-server checks to see if the client is authorized in step 520. If the client is 3 not authorized in step 520, then step 535 is executed. 4 5 In step 535, the semi-trusted web-server sends a 6 forbidden message to the client and returns to step 7 If the client is authorized in step 520, then step 525 is executed and the content is provided to the 8 client. Following step 525, the semi-trusted 9 web-server returns to step 505 and waits for more 10 11 messages.

Fig. 6 shows a second embodiment in which the client credential is presented to the semi-trusted web-server 600 by the origin web-server 610 and the origin server 610 redirects the client to the semi-trusted web-server 635 with HTTP redirection. This embodiment addresses the case when the origin web-server and semi-trusted web-server are not located in the same domain, although it can also be used if the servers are in the same domain. In this embodiment, the client makes a request to the origin web-server. The request is sometimes referred to the origin web-server by the semi-trusted The client is authenticated to the origin web-server. web-server with the authentication 630 component. At the origin web-server, the credential creation 620 component creates a client credential in the form of an HTTP cookie which includes client-specific environment information such as the apparent client Internet Protocol (IP) address and the HTTP header information. The presentation component 625 sends the client credential to the semi-trusted web-server 635, and then

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origin web-server 610 refers the client to the semi-trusted web-server by HTTP redirection including a reference to the client credential cookie in the HTTP redirection URL. When the client 605 makes a request to the semi-trusted web-server 635, the presentation component 615 gives a reference to the client credential cookie in the form of a URL query string, or if the client already has a client credential cookie the cookie is presented to the semi-trusted web-server. In the correlation component 635, the reference to the client credential cookie is used to select one of the stored cookies at the semi-trusted web-server, and then client-specific environment information such as the apparent client IP address and the HTTP header information is correlated to the client credential as illustrated in Fig. 12. If the client matches the client credential, then the semi-trusted web-server sets the cookie on the client for future requests and the semi-trusted web-server responds to the request using the access control of the user identifier specified in the cookie. The format and validation of the cookie are illustrated in Fig. 11.

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Fig. 7 shows a third embodiment in which correlation is aided by the origin web-server installing a first client-side program in the client browser that sends client-specific information to the origin web-server and by the semi-trusted web-server installing a second client-side program in the client browser that sends client-specific information to the semi-trusted web-server. This embodiment addresses the case when the origin web-server and semi-trusted web-server are

1 not located in the same domain, although it can also be 2 used if the servers are in the same domain. 3 embodiment further reduces the risk of using cookies for authentication by using more client-specific 5 environment information. In this embodiment, the 6 client makes a request to the origin web-server 7 (possibly being referred to the origin web-server by 8 the semi-trusted web-server). The client is 9 authenticated to the origin web-server with the 10 authentication 730 component, which may be enhanced 11 with the first client-side program. At the origin 12 web-server, the credential creation 720 component 413 creates a client credential in the form of an HTTP 114 cookie which includes client-specific environment 115 information such as the apparent client Internet **1**16 Protocol (IP) address and the HTTP header information, **47** as well as additional client-specific environment 18 information that is collected by the first client-side 19 program such as a hash of the local IP address and browser application process identifier. presentation component 725 sends the client credential 22 to the semi-trusted web-server 735, and then origin 23 web-server 710 refers the client to the semi-trusted 24 web-server by HTTP redirection including a reference to 25 the client credential cookie in the HTTP redirection 26 When the client 705 makes a request to the 27 semi-trusted web-server 735, the presentation component 715 gives a reference to the client credential cookie 29 in the form of a URL query string, or if the client 30 already has a client credential cookie the cookie is presented to the semi-trusted web-server. In the 32 correlation component 735, the reference to the client

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credential cookie is used to select one of the stored cookies at the semi-trusted web-server, and then client-specific environment information such as the apparent client IP address and the HTTP header information as well as additional client-specific environment information obtained by the correlation component 740 through the second client-side program is correlated to the client credential. If the client matches the client credential, then the semi-trusted web-server sets the cookie on the client for future requests and the semi-trusted web-server responds to the request using the access control of the user identifier specified in the cookie. The format and validation of the cookie are illustrated in Fig. 11.

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Fig. 8 shows the message exchange that occurs in the third embodiment as illustrated in Fig. 7. A client sends an HTTP request (800) to the origin web-server. The origin web-server replies with an authentication request message (805), which is presented to the user as a user ID and password prompt. The client replies with an authentication response message (810) including the authentication data. If the user ID and password are validated successfully at the origin web-server then the origin web-server sends a client program to the client (815). The client program runs on the client machine and gathers local environment information, e.g., process ID and user ID of the client's browser and local IP address. This information is hashed in the client for privacy and sent back as client specific correlation information to the origin web-server (820). The origin web-server creates a valid cookie as

1 described later in Fig. 11 and includes a reference to 2 the cookie in the redirection of the client (825). The 3 original web-server also stores the cookie at the semi-trusted web-servers (830) or in a directory 4 5 accessible to the semi-trusted web-servers. The client proceeds with sending the request to the semi-trusted 6 7 web-server (835) as indicated in the redirection. If 8 the semi-trusted web-server and original web-server 9 share one domain, then the client includes the cookie. 10 If the client does not include a cookie, then the 11 semi-trusted web-server looks up the cookie from where 12 it is stored by the origin web-server. If the 413 semi-trusted web-server does not receive a cookie from 114 the client and does not find a cookie in its accessible 415 storage, then the semi-trusted web-server redirects the 16 client to the original web-server and the process **4**17 restarts at message 800. The semi-trusted web-server =18 can implement protection against oscillating 119 redirections by allowing a client only a small number 20 of redirections in a specified time frame. If the 21 semi-trusted web-server has access to the client's cookie, then it sends a correlation client program to 22 23 the client (840), which gathers client environment 24 information (see 815) and send a hash of it back to the 25 proxy web-server (845). The semi-trusted web-server 26 then verifies whether the client's environment hash matches the environment hash stored in the cookie. If 27 28 not, then the semi-trusted web-server redirects the 29 client to the origin web-server and the scenario restarts (800). If it matches, then the semi-trusted 30 web-server validates the cookie (as described later in 31 32 Fig. 11). If the cookie is valid, then the semi-trusted

web-server extracts the client's credentials from the 1 2 cookie and continues with the authorization. If the 3 client's credentials are sufficient to access the requested information then access is provided to the 4 client, if not then the access is denied (850). The 5 6 semi-trusted web-server updates the cookie (as 7 described in Fig. 11), sends the updated cookie with message 850 to the client, and stores the updated 8 cookie in the accessible storage for semi-trusted 9 10 web-servers.

Fig. 9 shows a flowchart that describes the behavior of the origin web-server in the third embodiment as shown in Fig. 7. After starting the origin web-server (900), it waits for request messages (905). When receiving a client request, the origin web-server gathers client environment information (910) as described in Fig. 8 by sending a client side program to the client, which reports respective information. The origin web-server examines whether a valid cookie for this client exists (915). If yes, then the client is referred to a proxy web-server (920) and the origin web-server waits for other requests (905). Otherwise, the origin web-server authenticates the client (925) by sending an authentication request to the client and waiting for the authentication response (as described in Fig. 8, messages 805 and 810). If the authentication fails, then an access forbidden message is sent to the client (835) and the origin web-server waits for other requests (905). In this case, the origin web-server should audit the number of failed authentication procedures per client and lock accounts where a certain

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1 limit of successive authentication requests have 2 failed. If the authentication is successful and no 3 cookie is present then the origin web-server creates a cookie for the client as shown in Fig. 11, and sends 4 5 the cookie to the client together with a redirection message that refers the client to a semi-trusted 6 web-server (940). If origin and semi-trusted 7 8 web-servers do not share the same domain, the cookie 9 must be stored in a directory accessible to the 10 semi-trusted web-servers. Otherwise this is optional. 11 The origin web-server then returns to the state where 12 it waits for incoming requests (905).

> Fig. 10 shows a flowchart that describes the behavior of the semi-trusted web-server in the third embodiment as shown in Fig. 7. After starting the semi-trusted web-server (1000), it waits for incoming requests (1005). Next, the semi-trusted web-server checks whether it has as cookie for this client (1010). If the client does not submit a valid cookie with the request and the semi-trusted web-server does not have the client's cookie in other storage, e.g., in a directory, then the client is referred to the origin web-server (1015) for cookie creation. Otherwise, the semi-trusted web-server gathers client environment information (1020) using a client side program as described in Fig. 7. The semi-trusted web-server then verifies whether the gathered information and the client information stored in the cookie match (1025) as illustrated in Fig. 12. If it does not, then the client is referred to the origin web-server to create a valid cookie (1015). If it does, then the semi-trusted web-server updates

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1 the cookie (1035) as described in Fig. 13. Then the 2 semi-trusted web-server extracts the client's 3 credentials from the cookie and checks the 4 authorization of the client to access the requested 5 information (1040). If the client is not authorized to 6 access the requested information then the semi-trusted 7 web-server sends a forbidden message to the client 8 (1045) and returns to the waiting state (1005). If the 9 client is authorized then the client is provided access 10 (1050). Afterwards, the semi-trusted web-server returns 11 to the waiting state (1005).

> Fig. 11 illustrates an example of a client cookie that is created by the origin web-server. The cookie consists of two parts: one part is encrypted and one part is not encrypted. The encrypted part consists of the client's IP address as seen by the origin web-server (1100); optional client correlation information gathered either by a client-side program (1105); a hash of the client's request header as seen by the origin web-server (1110); the client's user identification as used for authorization by the origin and semi-trusted web-server (1115); optionally a random bit pattern B (1120); a time stamp including the creation time of the cookie (1125); a global time out value valid for the whole domain (1130) which is usually a fixed offset added to the creation time; and a cookie inactivity time-out (1135) which is a fixed offset added to the cookie creation time. The clear part of the cookie includes the key identification number (1140), which denotes the key used to encrypt the upper part of the cookie; the domain name of the

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origin web-server (1145); optionally a copy of the
encrypted random bit pattern B (1150); and finally a
digital signature (1155) over the fields above whereby
all fields are used in the clear for creating the
signature. Afterwards, the fields of the first part are
encrypted using the key Kc (1160) shared by the
semi-trusted web-server and the origin web-server.

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Fig. 12 shows a flowchart illustrating the process of validating a client cookie and returning client credentials in case the cookie is valid as part of the correlation procedures. The client cookie may be present in the HTTP request by the client, or the client may provide a reference to a cookie that is stored at the semi-trusted web-server or that is available to the semi-trusted web-server through an external device or directory. If no cookie is found, then the client is referred to the origin web-server so that the client can be authenticated and a cookie can be created. The cookie validation procedure starts at step 1200. In step 1205, the cookie is decrypted by using the domain identifier and key identifier to select an appropriate decryption key and by performing the decrypt operation. In step 1210, the global time-out and inactivity time-out fields are checked to see if the cookie has expired and the bit pattern is compared to the bit pattern copy. If the cookie has expired with either the global time-out or the inactivity time-out or if the bit pattern and bit pattern copy are not equal, then the cookie is invalid and step 1215 is executed. In step 1215, the process returns invalid and stops. If the cookie has not

expired and the bit patterns match in step 1210, then step 1220 is executed. In step 1220, if the key used to decrypt the cookie in step 1205 has been marked as compromised or if heightened security is desired, step 1225 is executed; otherwise step 1230 is executed. step 1225, the signature of the cookie is checked. Ιf the signature does not match, then the cookie is invalid and step 1215 is executed. If the signature does match, then step 1230 is executed. In step 1230, client-specific environment information is gathered by the semi-trusted web-server. Some client-specific environment information such as apparent IP address and HTTP header hash may come from the client connection, while other client-specific environment information such as a hash of the local user identifier, browser process identifier, and IP address may be sent to the semi-trusted web-server by the second client-side program. After step 1230, in step 1235 the hash of client-specific environment information is checked for equality with the client-specific environment information stored in the cookie. If there is not equality in step 1235, then the cookie is invalid and step 1215 is executed. If there is equality in step 1235, then step 1240 is executed. In step 1240, the client access credentials in the cookie are retrieved. In step 1245, the validation process reports valid and returns the client's credentials to the caller. These credentials are used throughout the authorization to decide whether to provide access or not to the client.

Fig. 13 shows a flowchart illustrating the process of updating a client cookie. After starting execution in

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step 1300, step 1305 is executed. Step 1305 checks to see whether or not the cookie has been marked valid in the validation process illustrated in Fig. 12. cookie is not valid in step 1305, then step 1330 is executed. In step 1330, the cookie is deleted from the client, local storage, and any cookie directory or storage device and step 1325 is executed. 1325, the cookie update process stops. If the cookie is valid in step 1305, then a new inactivity time-out is set in step 1310 and step 1315 is executed. 1315, the semi-trusted web-server checks the key management system to see if the key used by the client is still valid. The key may not be valid if the key has timed out or if the origin web-server has sent a message to the semi-trusted web-server indicating that the key is no longer valid. If in step 1315, the key is no longer valid, then step 1330 is executed. If in step 1315, the key is valid, then in step 1320 the cookie is encrypted with the same shared key indicated in the key identifier field and step 1325 is executed.

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Fig 14 illustrates the key entry used by semi-trusted web-servers and origin web-servers to protect part of the cookie both against unnoticed modification and disclosure. The key entry comprises of a key identification number (1400), the key itself (1405), and optionally a key time-out value (1410). At the beginning, the origin web-server creates a key entry with a new key, a new key identification number, and an empty time-out value and distributes it securely to all semi-trusted web-servers. The origin server includes in the appropriate cookie field (see Fig. 11), the key

identification of the key that is used when creating the cookie. Once a semi-trusted web-server is known to be compromised, the origin web-server issues a new key entry with a new key identification number to those semi-trusted web-servers that are not compromised; this new key issue triggers the semi-trusted web-servers to set a global time-out value to the old key (e.g. 5 seconds). After this time-out value expires, cookies using the old key are no longer accepted (see Fig. 13) and clients presenting them are redirected to the origin web-server for authentication and creation of a new cookie. This way, cookies and user information related to cookies that were known by the compromised web server lose their value for replay attacks.

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Fig. 15 shows a flowchart that illustrates the actions taken by the origin web-server when creating secure content for a client (1500). This part applies if the client retrieves information that shall be kept confidential even regarding the semi-trusted web-servers serving these data. Embedding secure content into information served by semi-trusted web-servers enhances scalability and security as most of the information retrieved is not highly sensitive and can be shared by many users leveraging caching in the semi-trusted web-server. The small amount of individual and sensitive data is retrieved transparently for the client and automatically by the semi-trusted web-server from the origin web-server. A special secure content handler, installed in the client, will present the secure content to the user in a way that makes secure and conventional data

distinguishable for the user. To enable secure content, the content provider determines the key Kclient, which it shares with the client (1505). It then encrypts the sensitive content with a key Kcontent (1510). The key Kcontent is encrypted with Kclient and added to the secure content (1515). The content is marked with a new tag (1520) and then served either directly to the client or to the semi-trusted web-server which embeds the secure content as part of the content it serves to this client (1525). The procedure ends with 1530. Using an independent key Kcontent to encrypt sensitive content supports the use of encrypted content for different users which do not share the key Kclient; the origin web-server just needs to encrypt the common key multiple times instead of the whole content.

Fig. 16 shows a flowchart that illustrates the actions taken by the semi-trusted web-server when it serves partly secure content (1600). If the secure content is retrieved by the semi-trusted web-server, then client identification information (e.g. the client's cookie) is determined (1605) and presented to the origin web-server when retrieving the secure content from the origin web-server (1610). This identification enables the origin web-server to determine the key Kclient (see Fig. 15). The semi-trusted web-server receives the secure content for the client from the origin web-server and serves the content to the client (1615). Whether secure content can be cached or not depends on the content itself and is determined by the origin web-server. The procedure serving the secure content part of a request ends with 1620.

1 Fig. 17 shows a flowchart that illustrates the actions 2 taken by the client when receiving secure content 3 (1700). First, the client's browser will check whether 4 it has a registered program for the type of content it 5 received (1705). If the client receives conventional 6 content then it processes the data with the existing 7 content handlers (1710) and returns (1715). If the 8 client receives data tagged as secure content, the 9 browser will automatically look for the respective [FF] registered program to handle this content (1720). The **#111** client recognizes secure content by its special content **1112** tag (see Fig. 18). If the browser does not have 1113 registered a handler for secure content, then the user 14 111 is prompted and must go through the installation 15 procedure. The client-side handler for secure content 16 is installed via a secure connection from the origin U17 web-server; throughout this process (1725), the secure key Kclient is installed in the client program as well. 18 19 Kclient is known only to the client-side program and 20 the content provider, e.g., a cgi script or servlet in 21 the origin web-server. Next, the client program as part 22 of the web-browser is triggered to handle the secure 23 content-type. The client-side program extracts the 24 protected key, Kp, attached to the secure content 25 (1730) and computes the decryption key Kcontent for the 26 secure content by decrypting the key Kp with its own 27 secret key Kclient (1735). The resulting key Kcontent 28 is used to decrypt the secure content and check the 29 integrity of the secure content (1740). The resulting 30 cleartext can be handled as usual by the web-browser or

- 1 can be presented by the client-side program in a way
- 2 that shows users which part of a page is secure content
- 3 and which part is not (1745).
- 4 Fig. 18 illustrates the new content-type that is
- attached to secure content (1800). A new tag (1805) 5
- 6 triggers the client's web-browser to activate the
- 7 client-side program (installed as described in Fig.
- 15). The protected key section (1810) comprises the 8
- 9 key, protected by Kclient, to decrypt the secure
- content. The tag is followed by the encrypted secure 10
- 11 content (1815). One possibility to implement the secure 12
 - content tag (1805) is to use a special file extension
- 113 (e.g., ".sec"), which is handled by a new content 14
 - handler. The file itself contains the protected content
- key (1810) and the secure content (1815). Another
- 15 116 possibility is to define a new HTML-Tag for secure
- 17 18 content, which implements the secure content tag (1805)
 - and a pointer to the file including the protected
- content key (1810) and the secure content (1815).
- 20 It is noted that the foregoing has outlined some of the
- 21 more pertinent objects and embodiments of the present
- 22 This invention may be used for many invention.
- 23 applications. Thus, although the description is made
- for particular arrangements and methods, the intent and 24
- concept of the invention is suitable and applicable to 25
- other arrangements and applications. It will be clear 26
- 27 to those skilled in the art that modifications to the
- disclosed embodiments can be effected without departing 28
- 29 from the spirit and scope of the invention.

- described embodiments ought to be construed to be
 merely illustrative of some of the more prominent
 features and applications of the invention. Other
 beneficial results can be realized by applying the
 disclosed invention in a different manner or modifying
- disclosed invention in a different manner or modifying the invention in ways known to those familiar with the
- 7 art.
- 8 The present invention can be realized in hardware, 9 software, or a combination of hardware and software. A 10 visualization tool according to the present invention **11** can be realized in a centralized fashion in one 12 computer system, or in a distributed fashion where **4**13 different elements are spread across several 1114 interconnected computer systems. Any kind of computer 15 system - or other apparatus adapted for carrying out **116** the methods described herein - is suitable. A typical 17 combination of hardware and software could be a general **U18** purpose computer system with a computer program that, 19 when being loaded and executed, controls the computer 20 system such that it carries out the methods described 21 herein. The present invention can also be embedded in a 22 computer program product, which comprises all the 23 features enabling the implementation of the methods 24 described herein, and which - when loaded in a computer 25 system - is able to carry out these methods.
 - Computer program means or computer program in the
 present context mean any expression, in any language,
 code or notation, of a set of instructions intended to
 cause a system having an information processing
 capability to perform a particular function either

- directly or after either or both of the following a)
- 2 conversion to another language, code or notation; b)
- 3 reproduction in a different material form.

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